# Week 9 Part I

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# Overview

- Dynamic allocation continued
- Heap versus stack
- Memory-related bugs
- Exam #2

# Dynamic Allocation

# Recall...

- Dynamic memory allocation allows us to request memory on the fly
- A way to use only what you need

```
// size_t is an integral defined
// elsewhere
void* malloc( size_t numBytes );
void* calloc( size_t num, size_t size );
void* realloc( void* ptr, size_t size );
```

# Recall...

Requesting some unknown number of integers at runtime

```
int numToAllocate;
scanf( "%i", &numToAllocate );
int* nums =
  malloc(sizeof( int ) * numToAllocate);
int nums2[ numToAllocate ]; // ERROR
```

# Multidimensional Arrays

- Need two separate allocations:
  - Array of columns
  - Each column individually

# Example

 Function that makes a matrix with a given number of rows and columns, all initialized to 0

```
int** makeMatrix( int rows, int cols ) {
  int** retval =
    calloc(rows, sizeof(int*));
  int row;
  for ( row = 0; row < rows; row++ ) {
    retval[row] =
      calloc(cols, sizeof(int));
  return retval;
```

# Question

• What differs here?

```
int** makeMatrix( int rows, int cols ) {
  int** retval =
    calloc(rows, sizeof(int*));
  int* temp = calloc(cols, sizeof(int));
  int row;
  for ( row = 0; row < rows; row++ ) {
    retval[ row ] = temp;
  return retval;
```

# With Structs

Works the same way

```
struct Foo {int x; int y;};
int main() {
   struct Foo* f =
      malloc( sizeof( struct Foo ) );
   f->x = 10;
   f->y = f->x;
   return 0;
}
```

# With Structs

Works the same way

```
struct Foo {int x; int y; };
int main() {
  int x;
  struct Foo** fs =
  calloc ( NUM STRUCTS,
           sizeof(struct Foo*);
  for (x = 0; x < NUM STRUCTS; x++) {
    fs[x] =
     malloc(sizeof(struct Foo));
  return 0;
```

# Example

- We want to calculate the sample standard deviation of some input numbers
- This formula is below:

$$s = \sqrt{\frac{1}{N-1} \sum_{i=1}^{N} (x_i - \overline{x})^2}$$

## Problems

- We need the average in order to calculate it, so we need to keep the numbers around
  - Can't simply calculate as we go
- Don't know how many numbers we need to read in
- Solution: Dynamic memory allocation!

stddev.c

# Heap vs. Stack

# Recall...

Say we have multiple function calls

```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

```
int main() {
 printf("%i",
          bar();
 return 0;
```

# Recall...

• How is this allocated in memory?

```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

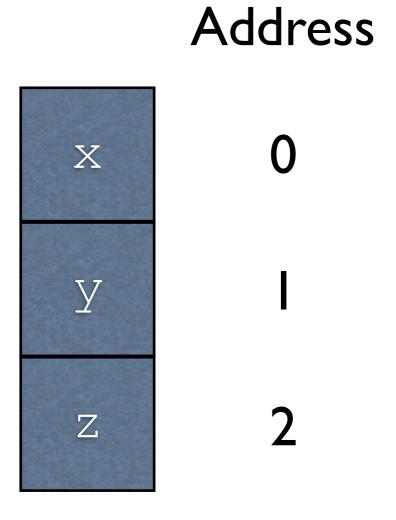
```
int main() {
 printf("%i",
          bar();
  return 0;
```

# One Possibility

```
int foo() {
                                   Address
  int x = 7;
  return x * 2;
                                X
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

# Pros

- Simple
- Fast



# Cons

- Wasteful (z is allocated but is unused in this code)
- Does not generally allow for recursion

```
int fact( int n ) {
  if ( n == 0 ) {
    return 1;
  } else {
    return n * fact(n-1);
  }
}
```

# Another Possibility

- Use a stack
  - As we call functions and declare variables, they get added to the stack
  - As function calls end and variables are no longer alive, they get removed from the stack
  - Do everything relative to the top of the stack

```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

```
int foo() {
  int x = 7;
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int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

Address



TOS

```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

#### Address



TOS

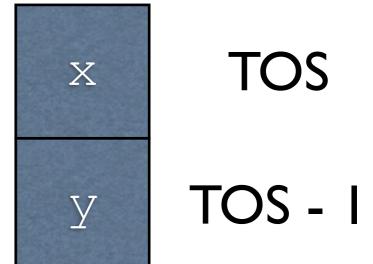
```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

#### Address



TOS

```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```



```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

# TOS Y TOS-I

```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```



```
int foo() {
  int x = 7;
  return x * 2;
int bar() {
  int y = 13;
  return y + foo();
int baz() {
  int z = 24;
  return z * 3;
```

# Notice

- The function baz was never called, and the variable z was thusly never put on the stack
- At all points, only exact as much memory as was needed was used

 Since stacks grow dynamically and allocate on the fly, we can have recursion

```
int fact (int n) {
  if (n == 0)
    return 1;
  } else {
    return n * fact(n-1);
int main() {
  fact(5);
  return 0;
```

```
int fact( int n ) {
  if (n == 0)
    return 1;
  } else {
   return n * fact(n-1);
int main() {
  fact(5);
  return 0;
```

```
fact (5)
n: 5
```

```
int fact( int n ) {
  if (n == 0)
    return 1;
  } else {
   return n * fact(n-1);
int main() {
  fact(5);
  return 0;
```

```
fact(4)
n: 4

TOS
fact(5)
n: 5
```

```
int fact( int n ) {
  if (n == 0)
    return 1;
  } else {
   return n * fact(n-1);
int main() {
  fact(5);
  return 0;
```

```
fact (3)
n: 3

fact (4)
n: 4

fact (5)
n: 5
```

```
int fact( int n ) {
  if (n == 0)
    return 1;
  } else {
   return n * fact(n-1);
int main() {
  fact(5);
  return 0;
```

```
fact(2)
fact(3)
 n: 3
fact (4)
 n: 4
fact (5)
 n: 5
```

#### Recursion

```
int fact( int n ) {
  if (n == 0)
    return 1;
  } else {
   return n * fact(n-1);
int main() {
  fact(5);
  return 0;
```

```
fact(1)
fact (2)
 n: 2
fact(3)
 n: 3
fact (4)
 n: 4
fact (5)
 n: 5
```

#### Recursion

```
int fact( int n ) {
  if (n == 0)
    return 1;
  } else {
   return n * fact(n-1);
int main() {
  fact(5);
  return 0;
```

```
fact (0)
fact(1)
 n: 1
fact(2)
 n: 2
fact(3)
 n: 3
fact (4)
 n: 4
fact (5)
 n: 5
```

### Stack Pros/Cons

- Only memory that is required is allocated
- Can have recursion
- Slower (everything now relative to the top of the stack instead of absolute)

## Question

• Is there anything odd with this code?

```
int* doSomething() {
  int x;
  return &x;
int main() {
  int w;
  int* p = doSomething();
  *p = 5;
  return 0;
```

```
int* doSomething() {
  int x;
  return &x;
int main() {
  int w;
  int* p = doSomething();
  *p = 5;
  return 0;
```



```
int* doSomething() {
  int x;
  return &x;
int main() {
  int w;
  int* p = doSomething();
  *p = 5;
  return 0;
```



```
int* doSomething() {
  int x;
  return &x;
int main() {
  int w;
  int* p = doSomething();
  *p = 5;
  return 0;
```



```
int* doSomething() {
  int x;
  return &x;
                                     TOS
                                X
int main() {
  int w;
                                      TOS - I
                                 W
  int* p = doSomething();
  *p = 5;
  return 0;
```

```
int* doSomething() {
  int x;
  return &x;
                                     TOS
                                X
int main() {
  int w;
                                W
  int* p = doSomething();
  *p = 5;
  return 0;
```

TOS - I

```
int* doSomething() {
  int x;
  return &x;
                                     TOS
                                W
int main() {
  int w;
  int* p = doSomething();
  *p = 5;
  return 0;
```

```
int* doSomething() {
  int x;
  return &x;
                                        TOS
                                   W
int main() {
  int w;
  int* p = doSomething();
  *p = 5;
                         Once TOS is updated, w's
  return 0;
                         address is the same as x's
                                address
```

### In General

- Stack allocation is done automatically
- Pointers to places on the stack are ok, as long as you make sure the address will always be alive
  - Pointers to the stack can be safely passed as function parameters, but not safely returned

## Heap

- Dynamic memory allocation allocates from a section of memory known as the heap
- Completely manual allocation
  - Allocate when you need it
  - Deallocate when you don't
- The computer assumes you know what you're doing

## Heap Allocation

Reconsider the code from before:

```
int* doSomething() {
// int x;
// return &x;
  return malloc( sizeof( int ) );
int main() {
  int w;
  int* p = doSomething();
  *p = 5;
  return 0;
```

## Heap Allocation

- This code is perfectly fine
  - There is nothing that will be automatically reclaimed
- Can pass pointers any which way

# Heap Allocation Bugs

## Memory Leak

- Allocated memory is not freed after it is needed (wasted)
- Generally, no pointers exist to the portion allocated, and so it can never be reclaimed
  - Why do we need a pointer?

```
void makeLeak() {
  malloc(sizeof(int));
}
```

# Dangling Pointer

- We keep a pointer to something that has been freed
- That memory can do just about anything after it is freed

```
int* dangle() {
  int* p = malloc( sizeof( int ) );
  free( p );
  *p = 5;
  return p;
}
```

#### Double Free

 We called free on something that was already freed

```
void doubleFree() {
  int* p = malloc( sizeof( int ) );
  free( p );
  free( p );
}
```

# On Memory Bugs

- Determining when something can be freed is generally difficult
  - Theoretically impossible to determine precisely in general
- Try to use stack allocation as much as possible

## Exam #2

#### Statistics

• A's: 8

Average: 78

• B's: 12

• Min: 52

• C's: 15

• Max: 97

• D's: 4

• F's: 2